

Translating the Cantor set by a random

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Goal:

Determine how much of the randomness in a “random” real can be cancelled by adding (or subtracting) a member of the Cantor set.

Further, determine the range of effective dimensions of points in a random translate of the Cantor set.

Let $E \subseteq \mathbf{R}^n$. The *diameter* of E , denoted $|E|$, is the maximum distance between any two points in E . We will use card for cardinality. A *cover* \mathcal{G} for a set E is a collection of sets whose union contains E , and \mathcal{G} is a δ -*mesh cover* if the diameter of each member G is at most δ . For a number $\beta \geq 0$, the β -*dimensional Hausdorff measure* of E , written $\mathcal{H}^\beta(E)$, is given by $\lim_{\delta \rightarrow 0} \mathcal{H}_\delta^\beta(E)$ where

$$\mathcal{H}_\delta^\beta(E) = \inf \left\{ \sum_{G \in \mathcal{G}} |G|^\beta : \mathcal{G} \text{ is a countable } \delta\text{-mesh cover of } E \right\}. \quad (1)$$

The *Hausdorff dimension* of a set E , denoted $\dim_{\text{H}}(E)$, is the unique number α where the α -dimensional Hausdorff measure of E transitions from being negligible to being infinitely large; if $\beta < \alpha$, then $\mathcal{H}^\beta(E) = \infty$ and if $\beta > \alpha$, then $\mathcal{H}^\beta(E) = 0$.

Let $S_\delta(E)$ denote the smallest number of sets of diameter at most δ which can cover E . The *upper box-counting dimension* of E is defined as

$$\overline{\dim}_B(E) = \limsup_{\delta \rightarrow 0} \frac{\log S_\delta(E)}{-\log \delta}.$$

For all E we have

$$\dim_H(E) \leq \overline{\dim}_B(E).$$

$$\dim_{\mathbb{H}}(A \times B) \geq \dim_{\mathbb{H}} A + \dim_{\mathbb{H}} B$$

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The *effective* (or *constructive*) β -dimensional Hausdorff measure of a set E , $c\mathcal{H}^\beta(E)$, is defined exactly in the same way as Hausdorff measure with the restriction that the covers be uniformly c.e. open sets. This yields the corresponding notion of the *effective* (or *constructive*) Hausdorff dimension of a set E , $\text{cdim}_H E$.

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Let

$$\mathcal{E}_{\leq \alpha} = \{x : \text{cdim}_{\text{H}}\{x\} \leq \alpha\}.$$

From the above, we know that the effective Hausdorff dimension of $\mathcal{E}_{\leq \alpha}$ satisfies $\text{cdim}_{\text{H}} \mathcal{E}_{\leq \alpha} = \alpha$; it turns out (Lutz) that also $\dim_{\text{H}} \mathcal{E}_{\leq \alpha} = \alpha$.

The *Kolmogorov complexity* of a string σ , denoted $K(\sigma)$, is the length (here we will measure length in ternary units) of the shortest program (under a fixed universal machine) which outputs σ . For a real number x , $x \upharpoonright n$ denotes the first n digits in a ternary expansion of x .

All sequences s of length n have $K(s) \leq n + O(\log n)$; most of them have $K(s) \geq n - O(1)$.

From work of Levin (\geq) and Mayordomo (\leq) we have for any real number x ,

$$\text{cdim}_{\text{H}}\{x\} = \liminf_{n \rightarrow \infty} \frac{K(x \upharpoonright n)}{n}. \quad (2)$$

We say a number is *Martin-Löf random* if it “passes” all Martin-Löf tests. A *Martin-Löf test* is a uniformly computably enumerable (c.e.) sequence of open sets $\{U_m\}_{m \in \mathbf{N}}$ with $\lambda(U_m) \leq 2^{-m}$, where λ denotes Lebesgue measure. A number x passes such a test if $x \notin \bigcap_m U_m$.

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Martin-Löf random reals have high initial segment complexity; indeed every Martin-Löf random real r satisfies $\lim_n K(r \upharpoonright n)/n = 1$.

Cantor set $\mathcal{C} \subseteq [0, 1]$:

$$\dim_{\text{H}} \mathcal{C} = \log_3 2 \approx .6309$$

$$\mathcal{C} + \mathcal{C} = [0, 2]$$

$$\frac{1}{2}\mathcal{C} + \frac{1}{2}\mathcal{C} = [0, 1]$$

$$E + C = [0, 2]$$

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$$\text{cdim}_H E = 1/2$$

$$B_1 = \{0, 1\}$$

$$B_2 = \{00, 02, 11\}$$

$$B_3 = \{000, 002, 021, 110, 112\}$$

$$B_4 = \{0000, 0002, 0011, 0200, 0202, 0211, 1100, 1102, 1111\}$$

$$B_5 = \{00000, 00002, 00021, 00112, 00210, 01221, 02012, \\ 02110, 02201, 10212, 11010, 11101, 11120, 11122\}$$

Lemma 1 (Lorentz). *There exists a constant c such that for any integer k , if $A \subseteq [0, k)$ is a set of integers with $|A| \geq \ell \geq 2$, then there exists a set of integers $B \subseteq (-k, k)$ such that $A + B \supseteq [0, k)$ with $|B| \leq ck \frac{\log \ell}{\ell}$.*

So we get E of constructive dimension $1 - \dim \mathcal{C}$ such that $E + \mathcal{C} = [0, 2]$.

Theorem 2. *Let $1 - \dim_{\text{H}} \mathcal{C} \leq \alpha \leq 1$ and let $r \in [0, 1]$. Then*

$$\dim_{\text{H}} [(\mathcal{C} + r) \cap \mathcal{E}_{\leq \alpha}] \geq \alpha - 1 + \dim_{\text{H}} \mathcal{C}.$$

Assume $1 - \dim_{\mathbb{H}} \mathcal{C} < \alpha < 1$. Split \mathbb{N} into $A \subseteq \mathbb{N}$ and \bar{A} ; then we can write $\mathcal{C} = \mathcal{C}_A + \mathcal{C}_{\bar{A}}$. Choose A of a suitable density $D = (1 - \alpha)/\dim_{\mathbb{H}} \mathcal{C}$ so that $\dim_{\mathbb{H}} \mathcal{C}_A$ comes out to be $1 - \alpha$. Then find a closed set E such that $\text{cdim}_{\mathbb{H}} E \leq \alpha$ and $\mathcal{C}_A + E = [0, 2]$. Let $F = 2 - E$, so that $F - \mathcal{C}_A = [0, 2]$. Let $S = \mathcal{C} \cap (F - r)$; it will suffice to show that $\dim_{\mathbb{H}} S \geq \alpha - 1 + \dim_{\mathbb{H}} \mathcal{C}$.

Now for each $z \in \mathcal{C}$ there exist unique points $v \in \mathcal{C}_A$ and $w \in \mathcal{C}_{\bar{A}}$ such that $v + w = z$; let p be the projection map which takes $z \in \mathcal{C}$ to its unique counterpart $w \in \mathcal{C}_{\bar{A}}$. For each $y \in \mathcal{C}_{\bar{A}}$ we have $r + y \in [0, 2] \subseteq F - \mathcal{C}_A$, so there exists $x \in \mathcal{C}_A$ such that $r + y \in F - x$, which gives $x + y \in S$ since $\mathcal{C}_A + \mathcal{C}_{\bar{A}} = \mathcal{C}$. Thus p maps S onto $\mathcal{C}_{\bar{A}}$. Since p is Lipschitz we have

$$\dim_{\mathbb{H}} S \geq \dim_{\mathbb{H}} \mathcal{C}_{\bar{A}} \geq \alpha - 1 + \dim_{\mathbb{H}} \mathcal{C}$$

because Lipschitz maps do not increase dimension. The theorem follows.

The set E is constructed as before, but the computation of $\text{cdim}_{\text{H}} E$ has an additional complication — we have not assumed that α is computable. It turns out that Kolmogorov complexity methods are useful here.

Specifically, let

$$A = \{ \lfloor y/D \rfloor : y \in \mathbf{N} \}.$$

Then initial segments of A are easy to describe:

$$K(A[n]) \leq 4 \log_3 n + O(1).$$

Now a straightforward computation shows that, if $\epsilon > 0$, then for all $x \in E$ and all sufficiently large k we have

$$K(x \upharpoonright m_k) \leq m_k[\alpha + \epsilon + o(1)],$$

which is enough to give $\text{cdim}_{\text{H}} E \leq \alpha$.

Theorem 3. *Let $1 - \dim_{\text{H}} \mathcal{C} \leq \alpha \leq 1$. For every Martin-Löf random real r ,*

$$\dim_{\text{H}} [(\mathcal{C} + r) \cap \mathcal{E}_{\leq \alpha}] \leq \alpha - 1 + \dim_{\text{H}} \mathcal{C}.$$

Theorem 4. *Let $1 - \dim_{\mathbb{H}} \mathcal{C} \leq \alpha \leq 1$ and let $r \in [0, 1]$ be Martin-Löf random. Then*

$$\dim_{\mathbb{H}} [(\mathcal{C} + r) \cap \mathcal{E}_{=\alpha}] = \alpha - 1 + \dim_{\mathbb{H}} \mathcal{C}.$$

Moreover,

$$\mathcal{H}^{\alpha-1+\dim_{\mathbb{H}} \mathcal{C}} [(\mathcal{C} + r) \cap \mathcal{E}_{=\alpha}] > 0.$$

Questions:

How much can the randomness of r be reduced by adding a Cantor set point if r was not completely random to begin with?

What about sets other than the Cantor set?